

## Error-Correcting Codes: Assignment #3

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**NOTE:** You can attempt a question after watching all video lectures up to and including the one listed in the question title.

1. **Combining two linear codes #1** (V3a) Let  $C_1$  be a binary  $(n_1, k_1, d_1)$ -code with generator matrix  $G_1$ . Let  $C_2$  be a binary  $(n_2, k_2, d_2)$ -code with generator matrix  $G_2$ . Let  $C$  be the binary code generated by the rows of the matrix

$$G = \left[ \begin{array}{c|c} G_1 & 0 \\ \hline 0 & G_2 \end{array} \right].$$

Determine the length  $n$ , dimension  $k$ , and distance  $d$  of  $C$ .

2. **Linear codes** (V3e)

Let  $C$  be an  $(n, k)$ -code over  $F = GF(q)$  with distance  $d$ .

- Prove that  $d \leq n - k + 1$ .
- Suppose that  $d = n - k + 1$ , and let  $S$  be any set of  $d$  coordinate positions. Prove that there is a codeword of weight  $d$  in  $C$  whose nonzero entries are in the coordinate positions of  $S$ . (For example, if  $n = 10$  and  $S = \{1, 3, 4, 7\}$ , then such a codeword might be  $c = (1, 0, 1, 1, 0, 0, 1, 0, 0, 0)$ .)
- Suppose that  $d = n - k + 1$ . Prove that the number of codewords of weight  $d$  in  $C$  is precisely  $(q - 1) \binom{n}{d}$ .

3. **Syndrome decoding** (V3f)

Consider the following parity check matrix  $H$  for a binary linear  $(n, k)$ -code  $C$  with distance  $d$ :

$$\begin{bmatrix} 1 & 1 & 1 & 1 & 0 & 1 & 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 & 1 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 1 & 1 & 0 & 1 & 0 & 1 & 1 \end{bmatrix}$$

- Determine  $n$ ,  $k$  and  $d$ .
- Set up a 1-1 correspondence between *syndromes* and *coset leaders*. Use the  $H$  given above.
- Decode each of the following vectors using the table of syndromes you constructed in (b):
  - $r_1 = (1010101010)$ .
  - $r_2 = (0011001100)$ .

4. **New codes from old ones** (V3f)

Let  $C$  be an  $(n, k, d)$ -binary code, and suppose that  $C$  has at least one codeword of odd weight. Let  $C'$  be the set of all even-weight codewords in  $C$ , and suppose that  $|C'| \geq 2$ .

- Prove that  $C'$  is a vector subspace of  $C$ .
- Prove that  $|C'| = \frac{1}{2}|C|$ .

Hint: Find a bijection between  $C'$  and the odd-weight codewords in  $C$ .
- Determine the length  $n'$  and dimension  $k'$  of  $C'$ .

(d) What can you say about the distance  $d'$  of  $C'$  in terms of  $d$ ?

5. **Golay codes** (V4b)

Decode each of the following received words using the decoding algorithm for  $C_{24}$  (the extended binary Golay code).

(a)  $r_1 = (0011\ 1000\ 0000\ 0100\ 1100\ 1110)$ .

(b)  $r_2 = (0000\ 0000\ 0011\ 1111\ 1101\ 1001)$ .

(c)  $r_3 = (1110\ 1000\ 0000\ 1001\ 0001\ 1101)$ .

(d)  $r_4 = (1111\ 0000\ 0000\ 0011\ 1010\ 0111)$ .